

# Buffer Occupancy of Statistical Multiplexers with Periodic Interchangeable Traffic in ATM Networks

Zhong Yu and Robert J. McEliece<sup>1</sup>

Dept. of Electrical Engineering, Caltech

Pasadena, CA 91125, U. S. A.

Email: zyu@systems.caltech.edu

**Abstract** — In this paper we analyze the buffer occupancy in a statistical multiplexer in ATM networks for a special type of traffic, namely, *periodic interchangeable* (PI) traffic. Certain generalized *Ballot theorem* is applied to analyze the problem. Explicit formulas for the expected buffer occupancy are derived.

## I. INTRODUCTION

An important concept in ATM networks is the efficient sharing of link capacities through statistical multiplexing of traffic sources. Buffering is provided at a statistical multiplexer to manage traffic fluctuations when the instantaneous rate of the aggregate incoming traffic exceeds the capacity of the outgoing link. The analysis of buffer occupancy for a statistical multiplexer with burst-constrained traffic is important to maintain quality-of-service guarantees in ATM networks. In this paper, we analyze the buffer occupancy for a statistical multiplexer with a special type of burst-constrained traffic, *periodic interchangeable* traffic.

## II. BALLOT THEOREM AND INTERCHANGEABLE RANDOM VARIABLES

**Ballot Theorem** [2]. If in a ballot candidate A scores  $a$  votes and candidate B scores  $b$  votes and if  $a \geq b$ , then the probability that throughout the counting the number of votes registered for A is always greater than the number of votes registered for B is given by

$$P(a, b) = \frac{a - b}{a + b}, \quad (1)$$

provided that all possible voting records are equally likely.

**Definition.** Random variables  $(X_1, X_2, \dots, X_M)$  are *cyclically interchangeable* if  $(X_1, X_2, \dots, X_M)$  take the value  $(x_1, x_2, \dots, x_M)$ , then they have the same probability for taking values of every cyclic permutation of  $(x_1, x_2, \dots, x_M)$ ; they are called *interchangeable* if they have the same probability for taking values of every permutation of  $(x_1, x_2, \dots, x_M)$ .

**Theorem (Takacs)** [2]. If  $(A_1, A_2, \dots, A_M)$  are cyclically interchangeable random variables such as  $\sum_{i=1}^M A_i = N$ , then

$$P\left\{\sum_{j=1}^i A_j < i, \forall i \in [1, M]\right\} = \begin{cases} 1 - \frac{N}{M} & \text{if } N < M \\ 0 & \text{otherwise} \end{cases} \quad (2)$$

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## III. STATISTICAL MULTIPLEXERS WITH PERIODIC INTERCHANGEABLE TRAFFIC

Consider a statistical multiplexer with an infinite buffer. Assume time is slotted and it takes a slot of time for the transmission line to transmit a cell. Suppose source  $i$ ,  $1 \leq i \leq K$ , generates  $n_i$  cells within a period of  $M$  slots; each source produces cells periodically with the same period  $M$ . The utilization factor  $\rho = \frac{N}{M} < 1$ . Define

$$\begin{cases} Q_k & = \text{buffer occupancy at the end of the } k\text{th slot} \\ A_k & = \text{number of arriving cells in the } k\text{th slot,} \end{cases}$$

and  $\bar{Q}$  the expected buffer occupancy. Note that  $\sum_{i=1}^M A_i = \sum_{i=1}^K n_i = N$ . Assume random variables  $A_1, A_2, \dots, A_M$  are interchangeable. Then, the traffic model is called *periodic interchangeable* (PI) traffic, which is a  $(\sigma, \rho)$  regulated traffic source with  $\sigma = \rho(M - 1)$ . A PI traffic in which each source generates cell uniformly within a period is called *uniform* PI (UPI) traffic. The UPI traffic in which each source generates only one cell within a period is called the *unit* UPI (UUPI) traffic [1]. Applying Takacs' generalized Ballot theorem, we can prove the following theorems:

**Theorem 1.** For a statistical multiplexer with UPI traffic,

$$\bar{Q} = \sum_{q=1}^N \sum_{r=1}^{N-q} \frac{M - N + q}{M - r} \sum_{\substack{\sum_{i=1}^K m_i = q + r \\ 0 \leq m_i \leq n_i}} \prod_{i=1}^K \frac{\binom{r}{m_i} \binom{M-r}{n_i - m_i}}{\binom{M}{n_i}}. \quad (3)$$

**Theorem 2.** For a statistical multiplexer with UUPI traffic, when  $K$  is sufficiently large,

$$\bar{Q} \begin{cases} = \frac{1}{2} \left( \sqrt{\frac{\pi K}{2}} - \frac{1}{3} \right) + o(1) & \rho = 1 \\ \leq \frac{\rho \alpha (1 - \rho)}{1 - \alpha} + o(1) & \rho < 1. \end{cases} \quad (4)$$

## IV. CONCLUSION

In this paper we applied Takacs' generalized ballot theorem to analyze the queue occupancy in a statistical multiplexer in ATM networks for a special type of burst-constrained traffic, *periodic interchangeable* (PI) traffic. Explicit formulas for the expected queue occupancy are derived for special PI traffic: UPI and UUPI. The result can shed light on the study of the worst case performance of a statistical multiplexer with burst-constrained traffic in ATM networks.

## REFERENCES

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