

# Construction of Asymptotically Good Low-Rate Error-Correcting Codes through Pseudo-Random Graphs

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A new technique, based on the pseudo-random properties of certain graphs, known as expanders, is used to obtain new simple explicit constructions of asymptotically good codes. In one of the constructions (construction  $\mathcal{C}_1$  below), the expanders are used to enhance Justesen codes by replicating, shuffling and then regrouping the code coordinates. For a given relative minimum distance  $\delta$  and alphabet size  $q$ , the codes thus obtained have rate  $R_1(\delta)$  which satisfies the inequality

$$R_1(\delta) \geq \gamma_0(1 - \delta) - \frac{\gamma_1}{\log_2 q}$$

for some positive constants  $\gamma_0$  and  $\gamma_1$ . In particular, for the low-rate neighborhood and for sufficiently large  $q$ , these codes lie above the Zyablov bound (cf. [3, Ch. 10]). For the small alphabet range (e.g.,  $GF(2)$ ), we obtain a second asymptotic good construction by using construction  $\mathcal{C}_1$  as an outer code in a concatenated scheme. The resulting codes, referred to as construction  $\mathcal{C}_2$ , have rate  $R_2(\delta)$  which is bounded from below by

$$R_2(\delta) \geq \max_{\delta \leq \mu \leq 1 - \frac{1}{q}} \gamma_0(1 - H_q(\mu)) \left(1 - \frac{\delta}{\mu}\right),$$

where  $H_q(x) \triangleq -x \cdot \log_q x - (1 - x) \cdot \log_q(1 - x) + x \cdot \log_q(q - 1)$ . Although construction  $\mathcal{C}_2$  lies below the Zyablov bound (the constant  $\gamma_0$  is approximately 0.021), they are still superior in the zero-rate neighborhood to previously-known explicit constructions, the best of which is due to Sugiyama et al. [4].

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These expander graphs, used in our construction, are  $\Delta$ -regular undirected graphs  $G = (V, E)$  such that, for any fixed real number  $\delta_0 \in (0, 1]$ , and for any subset of vertices  $B \subseteq V$  of size  $\geq \delta_0 |V|$ , the fraction of vertices in  $V$  which have at least one neighbor in  $B$  approaches unity "fast" as  $\Delta \rightarrow \infty$ . The specific expanders used are based on those introduced by Lubotzky et al. [1] and Margulis [2].

Given such a graph with  $n = |V|$  vertices and a finite field  $\Phi$ , we then define a mapping  $C_{\text{exp}} : \Phi^n \rightarrow (\Phi^\Delta)^n$ , such that every input  $n$ -tuple over  $\Phi$  of Hamming weight  $\geq \delta_0 n$  is mapped into an output  $n$ -tuple over  $\Phi^\Delta$  whose Hamming weight (measured over  $\Phi^\Delta$ ) is "close" to  $n$ . The above-mentioned construction  $\mathcal{C}_1$  is now obtained by applying the mapping  $C_{\text{exp}}$  to the codewords of Justesen codes, resulting in a code over the alphabet  $\Phi^\Delta$  whose rate is proportional to  $1/\Delta$  ( $\Delta$  and  $\Phi$  depend on the prescribed alphabet size  $q$  and relative minimum distance  $\delta$ ). Construction  $\mathcal{C}_2$  is obtained, in turn, by using Construction  $\mathcal{C}_1$  as an outer code, with each output symbol undergoing a second level of encoding by Wozencraft's ensemble codes.

## References

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